

Lecture/Exercise 2- First algorithms

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Deutsch-Jozsa algorithm

The Deutsch-Jozsa algorithm is historically important as showing a quantum advantage. It is also useful as a subroutine for other quantum algorithms, e.g., Grover's search.

Problem statement

Given $f : \{0, 1\}^n \rightarrow \{0, 1\}$, with $f(x) = \text{constant}$ ($f(x) = 1 \forall x$ or $f(x) = 0 \forall x$) or $f(x) = \text{balanced}$ (as many x give 0 as 1). $x = 0, \dots, 2^n - 1$ in binary language, i.e., $x = 00..0, 00..1, \dots, 11..1$.

Task of quantum/classical computer: determine if f is constant or balanced.

Available additional resource: a black box (an "Oracle machine") that implements f . So do not need to worry how efficient it is to call f , just how many times we need to call it.

Deterministic classical solution

In the worst case, a deterministic classical algorithm will need _____ evaluations of f .

Quantum algorithm

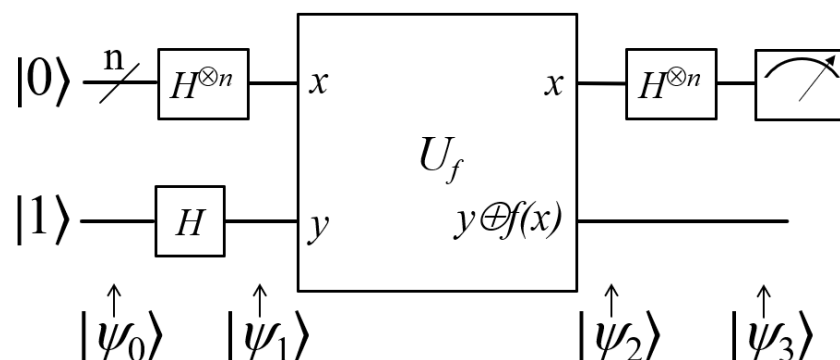


Figure 1: Circuit of the Deutsch-Jozsa algorithm.

0. Initial state: $|\psi_0\rangle = \text{_____}$ (first n are the register, last one is an "ancilla").

1. Hadamards: $|\psi_1\rangle = \frac{1}{\sqrt{2^{n+1}}} \sum_{x=0}^{2^n-1} |x\rangle \otimes (|0\rangle - |1\rangle)$
2. Call oracle: $|x\rangle|y\rangle \rightarrow |x\rangle|y \oplus f(x)\rangle$ (\oplus is addition modulo 2)

$$|\psi_2\rangle = \text{-----}$$

Rewrite, using

$$|f(x)\rangle - |1 \oplus f(x)\rangle = \text{-----} \text{ for } f(x) = 0, \text{ and}$$

$$|f(x)\rangle - |1 \oplus f(x)\rangle = \text{-----} \text{ for } f(x) = 1:$$

$$|\psi_2\rangle = \frac{1}{\sqrt{2^{n+1}}} \sum_{x=0}^{2^n-1} \text{-----} |x\rangle(|0\rangle - |1\rangle)$$

3. Apply another Hadamard to each qubit (from now on, ignore the last qubit): $|\psi_3\rangle = \frac{1}{2^n} \sum_{x=0}^{2^n-1} (-1)^{f(x)} \left[\sum_{y=0}^{2^n-1} (-1)^{x \cdot y} |y\rangle \right] = \frac{1}{2^n} \sum_{y=0}^{2^n-1} \left[\sum_{x=0}^{2^n-1} (-1)^{f(x)} (-1)^{x \cdot y} \right] |y\rangle$ where $x \cdot y = x_0 y_0 \oplus x_1 y_1 \oplus \dots \oplus x_{n-1} y_{n-1}$ is the sum of the bitwise product.

4. Probability of measuring $|0\rangle^{\otimes n}$:

$$|\langle 0|^{\otimes n} |\psi_3\rangle|^2 = \text{-----}$$

$$= \text{---} \text{ if } f(x) = \text{constant}$$

$$= \text{---} \text{ if } f(x) = \text{balanced}$$

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Which quantum features do you identify?

- -----
- -----
- -----
- -----
- -----

Discussion

The Deutsch-Jozsa algorithm

- produces an answer that is always correct requiring only **how many evaluations of f ?** -----
- yields an "oracle separation" between the classes ----- A common technique to make proves about complexity classes without having to know how a certain element can actually be realized.
- does not separate the classes ----- (the problem can be solved efficiently with probabilistic classical computer: to produce the correct answer with a high probability, k evaluations of the function suffice (failing with probability $\epsilon \leq 1/2^k$).